### CPE 460 Operating System Design Hey Process, Can we chat?

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### Hey Process! Can we communicate?

Processes executing concurrently in the operating system may be either independent processes or cooperating processes



A process is independent if it cannot affect or be affected by the other processes executing in the system.

Any process that does not share data with any other process is independent

### Cooperating processes need interprocess communication (IPC)

The operating system provides multiple mechanisms that allow processes to exchange data and information



# **Shared Memory**

An area of memory shared among the processes that wish to communicate







Kernel



Typically, a shared-memory region resides in the address space of the process creating the shared-memory segment.





Kernel



OS **prevents** one process from accessing another process's memory

Kernel

## **Shared Memory**

An area of memory shared among the processes that wish to communicate

Major issues is to provide mechanism that will allow the user processes to synchronize their actions when they access shared memory



## **Shared Memory**

An area of memory shared among the processes that wish to communicate

Major issues is to provide mechanism that will allow the user processes to synchronize their actions when they access shared memory



The consumer may have to wait for new items, but the producer can always produce new items.

### **POSIX Shared Memory**

POSIX shared memory is organized using **memory-mapped files**, which associate the region of shared memory with a file



https://github.com/torvalds/linux/blob/master/ipc/shm.c

https://www.safaribooksonline.com/library/view/linux-system-programming/0596009585/ch04s03.html

## **POSIX Shared Memory**

#include <stdio.h>
#include <stdib.h>
#include <string.h>
#include <fcntl.h>
#include <sys/shm.h>
#include <sys/stat.h>



int main()

/\* the size (in bytes) of shared memory object \*/
const int SIZE = 4096;
/\* name of the shared memory object \*/
const char \*name = "OS";
/\* strings written to shared memory \*/
const char \*message\_0 = "Hello";
const char \*message\_1 = "World!";

/\* shared memory file descriptor \*/
int shm\_fd;
/\* pointer to shared memory obect \*/
void \*ptr;

/\* create the shared memory object \*/
shm\_fd = shm\_open(name, O\_CREAT | O\_RDWR, 0666);

/\* configure the size of the shared memory object \*/
ftruncate(shm\_fd, SIZE);

/\* memory map the shared memory object \*/
ptr = mmap(0, SIZE, PROT\_WRITE, MAP\_SHARED, shm\_fd, 0);

/\* write to the shared memory object \*/
sprintf(ptr,"%s",message\_0);
ptr += strlen(message\_0);
sprintf(ptr,"%s",message\_1);
ptr += strlen(message\_1);

return 0;





#include <stdio.h>
#include <stdlib.h>
#include <fcntl.h>
#include <sys/shm.h>
#include <sys/stat.h>



int main()

/\* the size (in bytes) of shared memory object \*/
const int SIZE = 4096;
/\* name of the shared memory object \*/
const char \*name = "OS";
/\* shared memory file descriptor \*/
int shm\_fd;
/\* pointer to shared memory obect \*/
void \*ptr;

/\* open the shared memory object \*/
shm\_fd = shm\_open(name, O\_RDONLY, 0666);

/\* memory map the shared memory object \*/
ptr = mmap(0, SIZE, PROT\_READ, MAP\_SHARED, shm\_fd, 0);

/\* read from the shared memory object \*/
printf("%s",(char \*)ptr);

/\* remove the shared memory object \*/
shm\_unlink(name);

return 0;



### Message Passing

Communication takes place by means of messages exchanged between the cooperating processes

- 1. Establish a communication link between them
- 2. Exchange messages via send/receive



#### **Implementation issues:**

- How are communication links established?
- Can a link be associated with more than two processes?
- How many links can there be between every pair of communicating processes?
- What is the capacity of a link?
- Is a link unidirectional or bi-directional?

# Message Passing Communication

**Communication Link** 



#### **Direct Communication**

Processes must name each other explicitly send(, Normality, Normal

#### **Properties of communication link**

- Links are established automatically
- A link is associated with exactly one pair of communicating processes
- Between each pair there exists exactly one link
- The link may be unidirectional, but is usually bidirectional



Hard-Coding & Less Desirable

#### **Indirect Communication**

Messages are directed and received from mailboxes (ports)





#### **Properties of communication link**

- Link are established only if processes share a common mailbox
- A link may be associated with many processes
- Each pair of processes may share several communication links
- Link may be unidirectional or bi-directional



Flexible & More Desirable

### **Message Passing Operations**



# Message Passing Synchronization

**Blocking (Synchronous) Message Passing** 

send - the sender is blocked until the message is received receive - the receiver is blocked until a message is available

#### Non-blocking (Asynchronous) Message Passing

send - the sender sends the message and continue receive - the receiver receives: a valid message, or Null message

If both send and receive are blocking, we have a **rendezvous** 

# Message Passing Buffering



Zero Capacity no messages are queued on a link. Sender must wait for receiver (rendezvous)

Unbounded capacity infinite length. Sender never waits

Bounded Capacity Finite length messages Sender must wait if link full

### **POSIX Message Queues**

POSIX message queues is organized using virtual file system, and each message queue is pointed to by an mqueue\_inode\_info data structure





### **POSIX Message Queues**



https://www.softprayog.in/programming/interprocess-communication-using-posix-message-queues-in-linux

*Counts the number of files in a directory* 



A *pipe* acts as a conduit allowing *related processes* to communicate



Pipes

A conduit allowing *related processes* to communicate





Ordinary pipes is organized using **special file system** "not visible to user", and each pipe is pointed to by a pipe\_inode\_info data structure



https://github.com/torvalds/linux/blob/master/include/linux/fs.h#L553

https://github.com/torvalds/linux/blob/master/fs/pipe.c

https://books.google.com/books?id=LhQ7BAAAQBAJ&printsec=frontcover#v=onepage&q&f=true



https://www.softprayog.in/programming/interprocess-communication-using-posix-message-queues-in-linux

# Named Pipes "FIFO"

A named pipe acts as a bidirectional conduit allowing processes to communicate

No parent-child relationship required



http://en.wikipedia.org/wiki/Named\_pipe http://stackoverflow.com/questions/2784500/how-to-send-asimple-string-between-two-programs-using-pipes



